Modular Supervisory Synthesis for Unknown Plant Models Using Active Learning

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Modular Supervisors

Given a plant $G = \{G_1, G_2, \ldots, G_j\}$, let $K = \{K_1, K_2, \ldots, K_i\}$ be a set of automata describing the *specifications* of the system. A set of supervisors $S = \{S_1, S_2, \ldots, S_i\}$ can then be calculated for each specification K_i in such a way that the synchronous composition of the supervisors results in a maximally permissive *controllable* supervisor^a.

- Full synthesis
- Incremental synthesis

^aK. Åkesson, H. Flordal, and M. Fabian. "Exploiting modularity for synthesis and verification of supervisors". In: *IFAC Proceedings Volumes* 35.1 (2002). 15th IFAC World Congress, pp. 175–180. ISSN: 1474-6670.

Given a set of specification (automata) $K = \{K_1, K_2, ..., K_n\}$, a PSH $H = \langle M, E, S \rangle$ and a simulation, all corresponding to a system:

- Can we learn a supervisor?
- What properties will the supervisor satisfy: Non-blocking? Controllable? Maximally Permissive?

- One supervisor per specification
- To guarantee maximally permissive system each specification needs to be combined with modules defined in the PSH that
 - directly share uncontrollable events with the specification
 - indirectly share uncontrollable events with the specification

Definition (Event Dependence)

Given an alphabet Σ' and a set $M = m_1, m_2, \ldots, m_i$ of modules, let

 $\mathsf{Dep}(M, \Sigma') = \{m_i \in M \mid E(m_i) \cap \Sigma' \neq \emptyset\}$

Algorithm to define new modules

First initialize

$$\Sigma^{(1)} = \Sigma_u \cap \Sigma_{\mathcal{K}_i}$$
 $\mathcal{M}^{(1)}_{\mathcal{K}_i} = \mathsf{Dep}(\mathcal{M}, \Sigma^{(1)})$

Then repeat the following statements until $\Sigma^{(n+1)} = \Sigma^{(n)}$.

$$egin{aligned} \Sigma^{(n+1)} &= \Sigma^{(n)} \cup (\Sigma_{\mathcal{M}_{\mathcal{K}_i}^{(n)}} \cap \Sigma_u) \ && \mathcal{M}_{\mathcal{K}_i}^{(n+1)} &= \mathsf{Dep}(\mathcal{M}_{\mathcal{K}_i}^{(n)}, \Sigma^{(n)}) \end{aligned}$$

Defining supervisor modules (PSH + Spec)

- The supervisor learnt using the above PSH is maximally permissive.
- In some cases this could result in (inefficient) monolithic learning.

Incremental Learning and Synthesis

- Learn a supervisor for each specification and modules defined in the PSH that directly share uncontrollable events.
- Use existing tools to test if the obtained models are controllable, if not synthesize.
- For this, we need to learn additional modules: $M_G = M \setminus \bigcup_{K_i \in K} M_{K_i}$.

Updating E and S to handle specifications

$$E_{\mathcal{K}_i} = \Sigma_{\mathcal{K}_i} \cup \bigcup_{m \in \mathcal{M}_{\mathcal{K}_i}} E(m)$$
, and $S_{\mathcal{K}_i} = \bigcup_{m \in \mathcal{M}_{\mathcal{K}_i}} S(m)$.

Defining supervisor modules (PSH + Spec)

New PSH $\langle M', E', S', K' \rangle$

$$M' = M_K \cup M_G$$

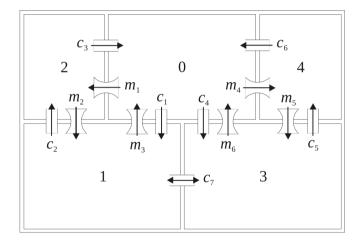
 $E'(m) = \begin{cases} E_{K_i}, & \text{if } m \in M_K \text{ and } m = M_{K_i} \\ E(m), & \text{if } m \in M_G \end{cases}$
 $S'(m) = \begin{cases} S_{K_i}, & \text{if } m \in M_K \text{ and } m = M_{K_i} \\ S(m), & \text{if } m \in M_G \end{cases}$
 $K'(m) = \begin{cases} K_i, & \text{where } m = M_{K_i} \\ undefined, & \text{otherwise.} \end{cases}$

We can now apply the modular learning (with slight modifications) to this new PSH to obtain the set of supervisors.

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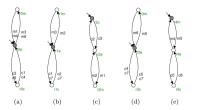


Fig. 2. Specifications for the different rooms $(Kr_0, Kr_1, Kr_2, Kr_3, Kr_4, \text{in that order})$ ensuring that only one of either the cat or the mouse can be present at a given time. Each state is identified using a unique name.

•
$$\Sigma_{Kr_0} = \{c_1, c_3, c_4, c_6, m_1, m_3, m_4, m_6\},\$$

•
$$\Sigma_{Kr_1} = \{c_1, c_2, c_7, m_2, m_3\},$$

•
$$\Sigma_{Kr_2} = \{c_2, c_3, m_1, m_2\},$$

•
$$\Sigma_{Kr_3} = \{c_4, c_5, c_7, m_5, m_6\},\$$

•
$$\Sigma_{Kr_4} = \{c_5, c_6, m_4, m_5\},\$$

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Cat and Mouse SUL/PSH

Simulation

- Let $\{R_0, R_1, R_2, R_3, R_4\}$ represent the different rooms.
- The simulator uses *var_c* and *var_m* to track the location of the animals.
- Initially, $var_c = R_2$ and $var_m = R_4$.

The PSH is defined as follows:

- *M* = {*Cat*, *Mouse*},
- $E(Cat) = \{c_1, c_2, c_3, c_4, c_5, c_6, c_7\},\$
- $E(Mouse) = \{m_1, m_2, m_3, m_4, m_5, m_6\},\$
- $S(Cat) = \{ var_c \},$
- $S(Mouse) = \{ var_m \},$

$$M_{Kr_1} = \{Cat\}, \text{ and } M_{Kr_3} = \{Cat\};$$

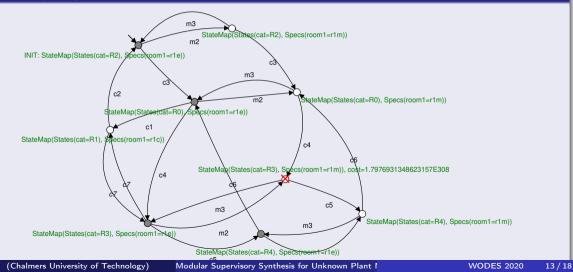
 $E_{Kr_1} = \Sigma_{Kr_1} \cup E(Cat), \text{ and } E_{Kr_3} = \Sigma_{Kr_3} \cup E(Cat);$
 $S_{Kr_1} = \{var_c, var_{Kr_1}\}, \text{ and } S_{Kr_3} = \{var_c, var_{Kr_3}\}.$

$$M_{K} = \{M_{Kr_{1}}, M_{Kr_{3}}\}, \text{ and } M_{G} = \{Mouse\};$$

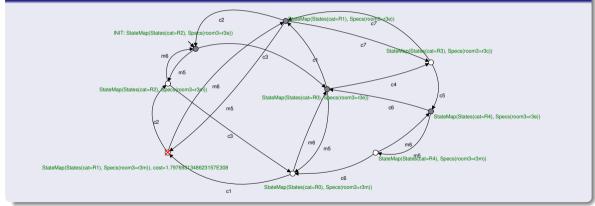
• The remaining specifications Kr_0 , Kr_2 , Kr_4 can be treated as supervisors.

- Explore the state space in a breath-first search manner using the PSH to do so modularly.
- When building supervisors:
 - For each transition identified check if the specification allows it.
 - Specifications can be tracked using state variables.
 - If the identified transition is uncontrollable:
 - Mark the originating state as forbidden and check for controllability for all coreachable states.
 - Forbidden states will not be explored further.

room 1 (M_{K1})



room 3 (M_{K3})

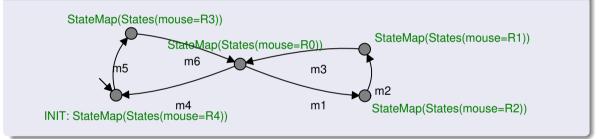


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Mouse



The resulting supervisors (and plants) can then be used in existing tools to generate a maximally permissive controllable and non-blocking supervisor.

- Not all decomposable systems are learnable.
 - Consider a simulation which has a state variable for each room $\{R_0, R_1, R_2, R_3, R_4\}$
 - These variable can take the value Cat or Mouse
 - Such a simulation is valid but cannot result in a modular model using this technique

Concluding remarks

- This demonstrates the possibility to obtain supervisors for existing systems when models are not present.
- Larger systems like the AGV example can be solved within a short time.
- PSH is tricky to define, and needs to be in conformance with the simulation; hence, knowledge of the simulation helps in creating a meaningful PSH.
- Accuracy and performance of this method depends upon the PSH.

Future Work

- Study the application of these techniques to a diverse range of examples.
- Learning richer formalism, like EFA.
- Simplify, and (maybe) automate the creation of the PSH.
- Learning different perspectives of modeling.

Thank you!

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